

Vježbe III

BCD kod (nastavak)

Oduzimanje u BCD kodu:

1. Umanjilac predstaviti komplementom devetke (jedinični komplement plus $10_{(10)}=1010_{(2)}$ koja se dodaje svakoj cifri pojedinačno, pri čemu se pretek sa cifre na cifru se odbacuje).
 2. Umanjenik sabrati sa komplementom devetke umanjioca (po pravilima sabiranja u BCD kodu).
 3. Pretek 1 označava pozitivan rezultat (predznak +, nema uticaja na brojčanu vrijednost) i dobijenom rezultatu se dodaje 1. Ukoliko preteka nema rezultat je negativan, pa je potrebno naći komplement devetke rezultata.

1. Izvršiti oduzimanje u BCD kodu:

- a) 495 i 278, b) 278 i 495.

a)

495		0100	1001	0101
- 278		0010	0111	1000
217	j.k.	1101	1000	0111
	d.k.	1010	1010	1010
		X 0111 X 0010 X 0001		
		0100	1001	0101
		0111	0010	0001

		1011	1011	0110
		0110	0110	
	①	0010	0001	0110
1 pretek \Rightarrow rezultat pozitiv			1	
		0010	0001	0111
		2	1	7 ₍₁₀₎

b)

278		0010	0111	1000
- 495		0100	1001	0101
- 217	j.k.	1011	0110	1010
	d.k.	1010	1010	1010
		X 0101 X 0000 X 0100		
		0010	0111	1000
		0101	0000	0100

		0111	0111	1100
				0110
	①	0111	1000	0010
0 pretek \Rightarrow rezultat negativan		1000	0111	1101
		1010	1010	1010
		X 0010 X 0001 X 0111		
		-2	1	7

Bulova algebra

$$\begin{array}{lll}
 A + 0 = A & A \times 0 = 0 & \\
 A + 1 = 1 & A \times 1 = A & \\
 A + A = A & A \times A = A & \\
 A + \bar{A} = 1 & A \times \bar{A} = 0 &
 \end{array}
 \quad \bar{\bar{A}} = A$$

De Morganova teorema

$$\begin{array}{l}
 \overline{A+B} = \bar{A} \times \bar{B} \\
 \overline{A \times B} = \bar{A} + \bar{B}
 \end{array}$$

2. Koristeći pravila Bulove agebre uprostiti izraze:

a) $A+\bar{A}B$ b) $A+B+\bar{A}\bar{B}$

a) $A+\bar{A}B = A(1+B)+\bar{A}B = A+AB+\bar{A}B = A+B(A+\bar{A}) = A+B$

b) $A+B+\bar{A}\bar{B} = A+B+\bar{A}+\bar{B} = 1$

3. Koristeći pravila Bulove agebre uprostiti izraze:

a) $F = \overline{AC} \times \overline{\bar{A}BC} \times \overline{\bar{B}C} + ABC\bar{C}$ b) $F = A\bar{B}\bar{C} + A\bar{B}C + ABC + A\bar{B}C$

a)

$$\begin{aligned}
 F &= \overline{AC} \times \overline{\bar{A}BC} \times \overline{\bar{B}C} + ABC\bar{C} = (\bar{A}+\bar{C}) \times (A+\bar{B}+\bar{C}) \times (B+\bar{C}) + ABC\bar{C} = \\
 &= (\bar{A}\bar{A} + \bar{A}\bar{B} + \bar{A}\bar{C} + A\bar{C} + \bar{C}\bar{B} + \bar{C}\bar{C}) \times (B+\bar{C}) + ABC\bar{C} = \\
 &= (\bar{A}\bar{B} + \bar{C}(\bar{A} + A + \bar{C} + \bar{B})) \times (B+\bar{C}) + ABC\bar{C} = \\
 &= (\bar{A}\bar{B} + \bar{C}(1 + \bar{C} + \bar{B})) \times (B+\bar{C}) + ABC\bar{C} = \\
 &= (\bar{A}\bar{B} + \bar{C}) \times (B+\bar{C}) + ABC\bar{C} = \\
 &= \bar{A}\bar{B}\bar{B} + \bar{A}\bar{B}C + \bar{C}B + \bar{C}\bar{C} + ABC\bar{C} = \bar{C}(\bar{A}\bar{B} + B + 1 + AB) = \bar{C}
 \end{aligned}$$

b)

$$\begin{aligned}
 F &= A\bar{B}\bar{C} + A\bar{B}C + ABC + A\bar{B}C = A\bar{B}(C + \bar{C}) + AB(C + \bar{C}) = \\
 &= A\bar{B} + AB = A(\bar{B} + B) = A
 \end{aligned}$$

4. Koristeći pravila Bulove agebre uprostiti funkciju.

$$\begin{aligned}
 F &= (A + \bar{C})(\bar{C} + EB) + (DC + EA)(\bar{B} + C) = (A + \bar{C})\bar{C}EB + (DC + EA)\bar{B}C = \\
 &= (A + \bar{C})\bar{C}(\bar{E} + \bar{B}) + \bar{B}CDC + \bar{B}CE\bar{A} = (A\bar{C} + \bar{C}\bar{C})(\bar{E} + \bar{B}) + \bar{B}CE\bar{A} = \\
 &= \bar{C}(A + 1)(\bar{E} + \bar{B}) + \bar{B}CE\bar{A} = \bar{C}(\bar{E} + \bar{B}) + \bar{B}CE\bar{A} = \bar{C}\bar{E} + \bar{C}B + \bar{B}CE\bar{A} = \\
 &= \bar{C}\bar{E} + \bar{B}C(1 + EA) = \bar{C}\bar{E} + \bar{B}C = \bar{C}(\bar{E} + \bar{B})
 \end{aligned}$$

5. Koristeći pravila Bulove ažube dokazati identitet $BC+ABD+A\bar{C}=BC+A\bar{C}$

$$\begin{aligned} BC+ABD+A\bar{C} &= BC+ABD(C+\bar{C})+A\bar{C} = BC+ACBD+A\bar{C}BD+A\bar{C} = \\ &= BC(1+AD)+A\bar{C}(BD+1) = BC+A\bar{C} \end{aligned}$$

6. Koristeći pravila Bulove algebre pokazati da važi $X_1 + X_2X_3 = (X_1 + X_2)(X_1 + X_3)$.

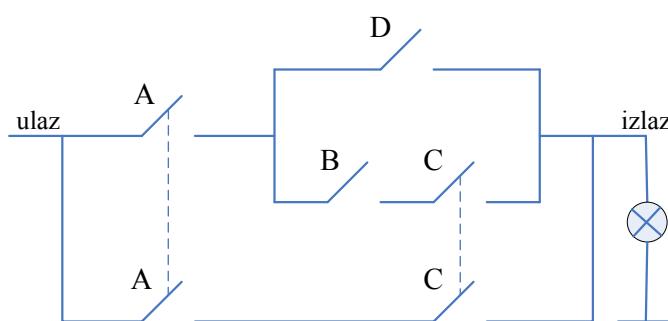
Posmatrajući desnu stranu uočavamo da nam nedostaju članovi X_1X_3 i X_1X_2 .

$$\begin{aligned} X_1 + X_2X_3 &= X_1(X_2 + \bar{X}_2) + X_2X_3 = \\ &= X_1X_2 + X_1\bar{X}_2 + X_2X_3 + X_1X_2 = \\ &= X_1(X_2 + \bar{X}_2) + X_2X_3 + X_1X_2 = \\ &= X_1 + X_2X_3 + X_1X_2 = \\ &= X_1(1 + X_3) + X_2X_3 + X_1X_2 = \\ &= X_1 + X_1X_3 + X_2X_3 + X_1X_2 = \\ &= X_1X_1 + X_1X_3 + X_2X_3 + X_1X_2 = \\ &= X_1(X_1 + X_3) + X_2(X_1 + X_3) = \\ &= (X_1 + X_2)(X_1 + X_3) \end{aligned}$$

7. Na slici je prikazana mreža sa relejima čiji kontakti A, B, C, D zauzimaju otvoren položaj u mirnom stanju.

a) Odrediti funkciju F koju realizuje data mreža uzimajući da je $F=1$ kada je put od ulaza ka izlazu mreže zatvoren.

b) Uprošćavanjem dobijene funkcije pokazati da njena vrijednost ne zavisi od položaja kontakta B.

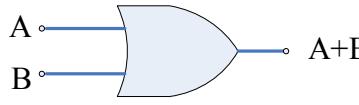
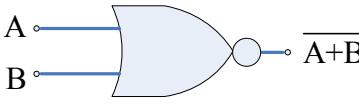
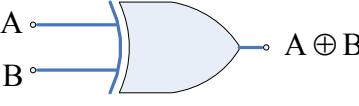
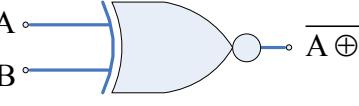
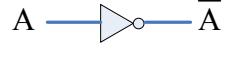


a) Donji put:
 $F=AC$ (i A i C moraju biti zatvoreni)

Gornji put:
 $F=A(D+BC)$

b) $F=AC+AD+ABC=AC(1+B)+AD=AC+AD=A(C+D) \Rightarrow$ ne zavisi od B

Logička kola i njihova funkcija

I kolo (množač)  $A \cdot B = AB$	<table border="1"> <thead> <tr> <th>A</th> <th>B</th> <th>AB</th> <th>\overline{AB}</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>0</td> <td>1</td> </tr> <tr> <td>0</td> <td>1</td> <td>0</td> <td>1</td> </tr> <tr> <td>1</td> <td>0</td> <td>0</td> <td>1</td> </tr> <tr> <td>1</td> <td>1</td> <td>1</td> <td>0</td> </tr> </tbody> </table>	A	B	AB	\overline{AB}	0	0	0	1	0	1	0	1	1	0	0	1	1	1	1	0	NI kolo  $A \cdot \overline{B} = \overline{AB}$
A	B	AB	\overline{AB}																			
0	0	0	1																			
0	1	0	1																			
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ILI kolo (sabirač)  $A + B = A+B$	<table border="1"> <thead> <tr> <th>A</th> <th>B</th> <th>$A+B$</th> <th>$\overline{A+B}$</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>0</td> <td>1</td> </tr> <tr> <td>0</td> <td>1</td> <td>1</td> <td>0</td> </tr> <tr> <td>1</td> <td>0</td> <td>1</td> <td>0</td> </tr> <tr> <td>1</td> <td>1</td> <td>1</td> <td>0</td> </tr> </tbody> </table>	A	B	$A+B$	$\overline{A+B}$	0	0	0	1	0	1	1	0	1	0	1	0	1	1	1	0	NILI kolo  $A + \overline{B} = \overline{A+B}$
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EKSCLUZIVNO ILI kolo  $A \oplus B = A \oplus B$	<table border="1"> <thead> <tr> <th>A</th> <th>B</th> <th>$A \oplus B$</th> <th>$\overline{A \oplus B}$</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>0</td> <td>1</td> </tr> <tr> <td>0</td> <td>1</td> <td>1</td> <td>0</td> </tr> <tr> <td>1</td> <td>0</td> <td>1</td> <td>0</td> </tr> <tr> <td>1</td> <td>1</td> <td>0</td> <td>1</td> </tr> </tbody> </table>	A	B	$A \oplus B$	$\overline{A \oplus B}$	0	0	0	1	0	1	1	0	1	0	1	0	1	1	0	1	EKSCLUZIVNO NILI kolo  $A + \overline{B} = \overline{A \oplus B}$
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0	0	0	1																			
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INVERTOR  $A \rightarrow \overline{A}$	<table border="1"> <thead> <tr> <th>A</th> <th>\overline{A}</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>1</td> </tr> <tr> <td>1</td> <td>0</td> </tr> </tbody> </table>	A	\overline{A}	0	1	1	0															
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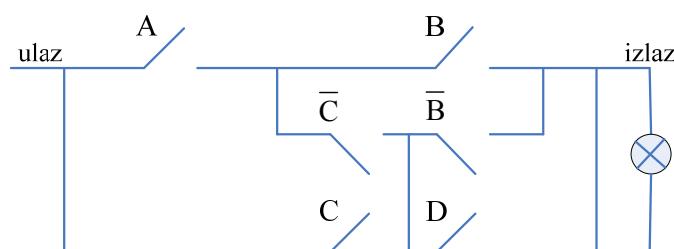
*Napomena:

Dodavanjem invertora prije ulaska signala u posmatrano kolo invertuje se samo taj signal!

Na primjer:



8. Na slici je prikazana mreža sa položajima kontakta releja u mirnom stanju, pri čemu je kontakt zatvoren ako je vrijednost promjenljive koja njima upravlja jednak logičkoj jedinici. Odrediti funkciju spoljnih puteva ulaz-izlaz, i po mogućnosti uprostiti datu mrežu, a zatim je predstaviti pomoću logičkih kola.



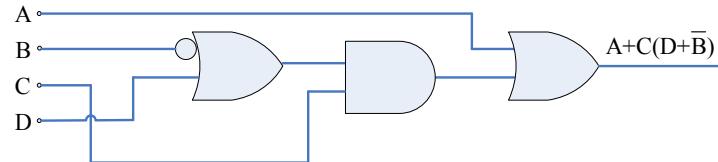
$$F = AB + A\bar{C}\bar{B} + CD + C\bar{B} + C\bar{C}B + A\bar{C}D = AB + A\bar{C}(\bar{B} + D) + C(D + \bar{B}) = AB + (\bar{B} + D)(A\bar{C} + C) =$$

pri čemu je:

$$A\bar{C} + C = A\bar{C} + C(A + \bar{A}) = A\bar{C} + AC + \bar{A}C + AC = A(C + \bar{C}) + C(A + \bar{A}) = A + C$$

pa dobijamo da je funkcija:

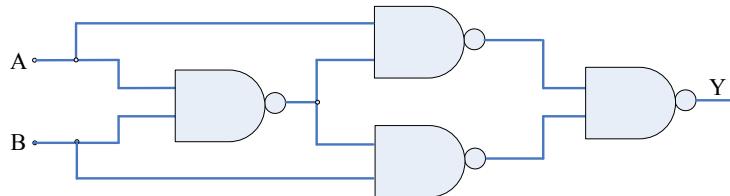
$$AB + (\bar{B} + D)(A + C) = AB + A\bar{B} + AD + C(\bar{B} + D) = A(B + \bar{B}) + AD + C(\bar{B} + D) = A(1 + D) + C(\bar{B} + D) = A + C(D + \bar{B})$$



9. Formirati „ekskluzivno ILI“ kolo koristeći samo NI kola.

$$Y = A \oplus B = \overline{AB} + A\bar{B}$$

$$Y = \overline{AB} + A\bar{B} + \overline{AA} + \overline{BB} = A(\overline{A} + \overline{B}) + B(\overline{A} + \overline{B}) = \overline{A}\overline{B} + B\overline{A} = \overline{\overline{A}\overline{B}} + \overline{B\overline{A}} = (\overline{A}\overline{B})(B\overline{A})$$



10 Realizovati funkciju $Y = \overline{AB} + AB$ koristeći samo NILI kola.

$$Y = \overline{AB} + AB = \overline{A + B} + \overline{AB} = \overline{A + B} + \overline{\overline{A} + \overline{B}} = \overline{\overline{A} + \overline{B} + \overline{A} + \overline{B}}$$

Ukoliko na oba ulaza NILI kola dovedemo isti signal dobijamo invertor:



što nam je potrebno za krajnje komplementiranje u posmatranom izrazu. Primjetimo da konstruisanje invertora na isti način moguće učiniti upotrebom NI kola.

